# Nonlinear Arithmetic in PVS 

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## Nonlinear Arithmetic

Interval can solve problems like

```
ex_ba : LEMMA
    x ## [|-1/2,0|] IMPLIES
    abs(ln}(1+x) - x) - epsilon <= 2*sq(x
```

Bernstein can solve problems like:

```
p1 : LEMMA
```

    FORALL ( \(\mathrm{x}, \mathrm{y}: \mathrm{real}\) ) : \(-0.5<=\mathrm{x}\) AND \(\mathrm{x}<=1\) AND
                            \(-2<=y\) AND \(y<=1\) IMPLIES
        \(4 * x^{\wedge} 2-(21 / 10) * x^{\wedge} 4+(1 / 3) * x^{\wedge} 6+(\mathrm{x}-3) * \mathrm{y}-4 * \mathrm{y}^{\wedge} 2+4 * \mathrm{y}^{\wedge} 4>-3.4\)
    p2 : LEMMA
    EXISTS ( \(\mathrm{x}, \mathrm{y}: r \mathrm{real}\) ) : \(-0.5<=\mathrm{x}\) AND \(\mathrm{x}<=1\) AND
                            \(-2<=y\) AND \(y<=1\) AND
        \(4 * x^{\wedge} 2-(21 / 10) * x^{\wedge} 4+(1 / 3) * x^{\wedge} 6+(x-3) * y-4 * y^{\wedge} 2+4 * y^{\wedge} 4<-3.39\)
    These lemmas are proved by executing a single command!

## Interval

http://shemesh.larc.nasa.gov/people/cam/Interval

- Interval is a PVS package for interval analysis.
- The package consists of:
- The library interval_arith, which presents a formalization of interval analysis for real-valued functions including: trigonometric functions, logarithm and exponential functions, square root, absolute value, etc.
- The strategy numerical, which implements a provably correct branch-and-bound interval analysis algorithm.
- Interval is part of the NASA PVS Libraries.


## A Simple Problem

Prove that the turn rate of an aircraft with a bank angle of $35^{\circ}$ is greater than $3^{\circ}$ per second.


Prove that the turn rate of an aircraft with a bank angle of $35^{\circ}$ is greater than $3^{\circ}$ per second.

```
IMPORTING interval_arith@strategies
g:posreal=9.8 %[m/s^2]
v:posreal=250*0.514 %[m/s]
tr(phi:(Tan?)): MACRO real = g*tan(phi)/v
tr_35 : LEMMA
    3*pi/180 <= tr(35*pi/180)
```

```
    tr_35 :
    |-------
{1} 3 * pi / 180 <= g * tan(35 * pi / 180) / v
Rule? (numerical)
```

Evaluating formula using numerical approximations,
Q.E.D.

Note that pi is the mathematical irrational number $\pi$ and $\tan$ is the trigonometric function tan.

```
tr_35 :
    |-------
{1} 3* pi / 180<= g * tan(35* pi / 180) / v
```

Rule? (numerical)
Evaluating formula using numerical approximations,
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Note that pi is the mathematical irrational number $\pi$ and $\tan$ is the trigonometric function tan.

## A Simple Property of Logarithms

```
G(x:real|x < 1): MACRO real = 3*x/2 - ln(1-x)
A_and_S : LEMMA
    let x = 0.5828 in
        G(x) > 0
```

A Simple Property of Logarithms

```
A_and_S :
    |-------
{1} LET x = 0.5828 IN 3 * x / 2 - ln(1 - x) > 0
Rule? (numerical)
Evaluating formula using numerical approximations,
Q.E.D.
```

Note that $\ln$ is natural logarithm function.

\{1\} LET $\mathrm{x}=0.5828$ IN $3 * \mathrm{x} / 2-\ln (1-\mathrm{x})>0$

Rule? (numerical)
Evaluating formula using numerical approximations, Q.E.D.

Note that ln is natural logarithm function.

```
{-1} x ## [| 0, 2 |]
    |-------
{1} sqrt(x) + sqrt(3) < pi + 0.1
```

Rule? (numerical :vars "x")
Evaluating formula using numerical approximations,
Q.E.D.

## Interval Arithmetic

```
{-1} x ## [| 0, 2 |]
```

|-------
\{1\} $\operatorname{sqrt}(x)+\operatorname{sqrt}(3)<\mathrm{pi}+0.1$
Rule? (numerical :vars "x")
Evaluating formula using numerical approximations, Q.E.D.

Prove that for all $x \in\left[-\frac{1}{2}, 0\right]$,

$$
|\ln (1+x)-x|-\epsilon \leq 2 x^{2},
$$

```
where }\epsilon=0.15:\mp@subsup{}{}{1
ex_ba : LEMMA
            x ## [|-1/2,0|] IMPLIES
        abs(ln(1+x) - x) - epsilon <= 2*sq(x)
```

${ }^{1}$ Thanks to Behzad Akbarpour.
instint
ex_ba :
|-------
\{1\} FORALL (x: real):
x \#\# $[|-1 / 2,0|]$ IMPLIES $\operatorname{abs}(\ln (1+\mathrm{x})-\mathrm{x})-0.15<=2 * \mathrm{sq}(\mathrm{x})$
Rule? (skeep)
ex_ba
\{-1\} x \#\# [| -1 / 2, 0 |]
\{1\} $\operatorname{abs}(\ln (1+x)-x)-0.15<=2 * s q(x)$

Rule? (numerical :vars (("x" 10)))
Evaluating formula using numerical approximations,
Q.E.D.

```
ex_ba :
    |-------
{1} FORALL (x: real):
    x ## [|-1/2,0|] IMPLIES abs(ln(1+x)-x)-0.15 <= 2*sq(x)
Rule? (skeep)
ex_ba :
{-1} x ## [l -1 / 2, 0 |]
    |-------
{1} abs(ln}(1+x)-x)-0.15<= 2* sq(x
Rule? (numerical :vars (("x" 10)))
Evaluating formula using numerical approximations,
Q.E.D.
```

ex_ba :
|-------
\{1\} FORALL (x: real):
x \#\# [|-1/2,0|] IMPLIES abs $(\ln (1+x)-x)-0.15<=2 * s q(x)$
Rule? (skeep)
ex_ba :
\{-1\} x \#\# [| -1 / 2, 0 |]
|-------
$\{1\} \operatorname{abs}(\ln (1+x)-x)-0.15<=2 * \operatorname{sq}(x)$

Rule? (numerical :vars (("x" 10)))
Evaluating formula using numerical approximations, Q.E.D.

## Bernstein

- Bernstein is a PVS package for solving multivariate polynomial global optimization problems using Bernstein polynomials.
- The package consists of:
- The library Bernstein, which presents a formalization of an efficient representation of multivariate polynomials.
- The strategy bernstein, which discharges simply quantified multivariate polynomial inequalities on closed/open ranges.
- Grizzly, which is a prototype client-server tool for solving global optimization problems.
- Bernstein is part of the NASA PVS Libraries.


## Solving Polynomial Inequalities

IMPORTING Bernstein@strategy
p1 : LEMMA
FORALL (x,y:real): -0.5 <= x AND $x<=1$ AND
-2 <= y AND y <= 1 IMPLIES $4 * x \wedge 2-(21 / 10) * x \wedge 4+(1 / 3) * x^{\wedge} 6+(x-3) * y-4 * y^{\wedge} 2+4 * y^{\wedge} 4>-3.4$
p2 : LEMMA
EXISTS (x,y:real): -0.5 <= x AND $x$ <= 1 AND
-2 <= y AND y <= 1 AND $4 * x^{\wedge} 2-(21 / 10) * x^{\wedge} 4+(1 / 3) * x^{\wedge} 6+(x-3) * y-4 * y^{\wedge} 2+4 * y^{\wedge} 4<-3.39$
|-------
\{1\} FORALL (x, y: real):

$$
\begin{aligned}
& -0.5<=\mathrm{x} \text { AND } \mathrm{x}<=1 \text { AND }-2<=\mathrm{y} \text { AND } \mathrm{y}<=1 \text { IMPLIES } \\
& 4 * \mathrm{x}^{\wedge} 2-(21 / 10) * \mathrm{x}^{\wedge} 4+(1 / 3) * \mathrm{x}^{\wedge} 6+(\mathrm{x}-3) * \mathrm{y}-4 * \mathrm{y}^{\wedge} 2+4 * \mathrm{y}^{\wedge} 4>-3.4
\end{aligned}
$$

Rule? (bernstein)
Proving polynomial inequality using Bernstein'basis,
Q.E.D.

\{1\} FORALL (x, y: real):

$$
\begin{aligned}
& -0.5<=\mathrm{x} \text { AND } \mathrm{x}<=1 \text { AND }-2<=\mathrm{y} \text { AND } \mathrm{y}<=1 \text { IMPLIES } \\
& 4 * \mathrm{x}^{\wedge} 2-(21 / 10) * \mathrm{x}^{\wedge} 4+(1 / 3) * \mathrm{x}^{\wedge} 6+(\mathrm{x}-3) * \mathrm{y}-4 * \mathrm{y}^{\wedge} 2+4 * \mathrm{y}^{\wedge} 4>-3.4
\end{aligned}
$$

Rule? (bernstein)
Proving polynomial inequality using Bernstein'basis, Q.E.D.
|-------
\{1\} EXISTS (x, y: real):

$$
\begin{gathered}
-0.5<=\mathrm{x} \text { AND } \mathrm{x}<=1 \text { AND }-2<=\mathrm{y} \text { AND } \mathrm{y}<=1 \text { AND } \\
4 * \mathrm{x}^{\wedge} 2-(21 / 10) * \mathrm{x}^{\wedge} 4+(1 / 3) * \mathrm{x}^{\wedge} 6+(\mathrm{x}-3) * \mathrm{y}-4 * \mathrm{y}^{\wedge} 2+4 * \mathrm{y}^{\wedge} 4<-3.39
\end{gathered}
$$

Rule? (bernstein)
Proving polynomial inequality using Bernstein's basis,
Q.E.D.
|-------
\{1\} EXISTS (x, y: real):

$$
\begin{gathered}
-0.5<=\mathrm{x} \text { AND } \mathrm{x}<=1 \text { AND }-2<=\mathrm{y} \text { AND } \mathrm{y}<=1 \text { AND } \\
4 * \mathrm{x}^{\wedge} 2-(21 / 10) * \mathrm{x}^{\wedge} 4+(1 / 3) * \mathrm{x}^{\wedge} 6+(\mathrm{x}-3) * \mathrm{y}-4 * \mathrm{y}^{\wedge} 2+4 * \mathrm{y}^{\wedge} 4<-3.39
\end{gathered}
$$

Rule? (bernstein)
Proving polynomial inequality using Bernstein's basis, Q.E.D.

## Reflection

Both Interval and Bernstein use computation reflection


Both try to prove the result on a large box:


## Reflection

Interval and Bernstein each have a function that can (sometimes) tell whether the result holds on a particular box.


If that function returns unknown, then the box is split in two:


## Reflection

The two halves of the big box are now considered separately


Perhaps we can prove it on the right but not the left sub-box:


## Reflection



This turns the proof tree into


## Reflection

Now we split the left hand box into two smaller pieces:


Perhaps the result can be be proved on each of these boxes:



This turns the proof tree into


Proof Splits
Proof Terminates

## Reflection

Sometimes the proof tree can get very large:


With 100s of splits, the proof is infeasible in the PVS prover language

## Reflection

Instead of having PVS develop a proof of that looks like


There is a recursive reflection function in PVS whose execution looks like


- The proof tree happens entirely in LISP
- All of the proofs have the same length in PVS, for Interval and Bernstein
- Complicated problems could not be solved in PVS without using computational reflection in this way.


## Reflection

There is a generic function in the structures library that defines a recursive splitting algorithm for arbitrary types.


It can be used to solve almost any binary branching problem

## Reflection

```
branch_and_bound(simplify,evaluate,branch, subdivide,denorm, combine,prune,le,ge, select,accumulate,maxdepth)
                    (obj,dom,acc,(dirvars|length(dirvars) <= maxdepth)) :
RECURSIVE Output =
LET nobj = simplify(obj),
        thisans = evaluate(dirvars,dom,nobj),
        newacc1 = IF none?(acc) THEN thisans ELSE accumulate(TRUE,some(acc),thisans) ENDIF,
        thisout = mk_out(thisans,ge(dirvars,newacc1,thisans),length(dirvars),0)
    IN
        IF length(dirvars)=maxdepth OR le(thisans) OR thisout`exit OR prune(dirvars,newacc1,thisans) THEN
        thisout
    ELSE
        (dir,v) = select(dirvars,newacc1, dom,nobj),
        funsplit = branch(v,nobj),
        (sp1,sp2) = IF dir THEN (funsplit`1,funsplit`2) ELSE (funsplit`2,funsplit`1) ENDIF,
        (dom1,dom2) = IF dir THEN (domsplit`1,domsplit`2) ELSE (domsplit`2,domsplit`1) ENDIF,
        firstout = branch_and_bound(simplify,evaluate,branch, subdivide,denorm,combine,
                    sp1,dom1,Some(newacc1), pushDirVar((dir,v),dirvars))
        IN
        IF firstout`exit THEN
            mk_out(combine(v,denorm((dir,v),firstout`ans),thisans),
            TRUE,firstout`depth,firstout`splits+1)
            ELSE
                newacc2 = accumulate(FALSE,newacc1,firstout`ans),
                secondout = branch_and_bound(simplify,evaluate,branch,subdivide,denorm,combine,
                    e, select, accumulate, maxdepth)
                                    (sp2,dom2, Some(newacc2), pushDirVar((NOT dir,v), dirvars))
            (real1,real2) = IF dir THEN (firstout,secondout) ELSE (secondout,firstout) ENDIF
            IN
                    secondout`exit,
                        max(firstout`depth, secondout`depth),
                        firstout`splits+secondout`splits+1)
            ENDIF
ENDIF
MEASURE maxdepth-length(dirvars)
```

- This algorithm can be evaluated by (eval-formula)
- ... and therefore, it can be used for computational reflection
- ... as long as everything it has to compute is a ground term


## Reflection



Yogi Berra: "It aint over 'til it's over"

Interval and Bernstein are not perfect
This algorithm may not terminate, even with Interval and Bernstein
There are some inequalites that are true that will not prove in a reasonable amount of time

THE END

THE END

